

# Theoretical and Empirical Analyses of the Effectiveness of Metamorphic Relation Composition

Kun Qiu, Zheng Zheng\*, *Senior Member, IEEE*, Tsong Yueh Chen, *Member, IEEE*, and Pak-Lok Poon, *Member, IEEE*

**Abstract**—Metamorphic Relations (MRs) play a key role in determining the fault detection capability of Metamorphic Testing (MT). As human judgement is required for MR identification, systematic MR generation has long been an important research area in MT. Additionally, due to the extra program executions required for follow-up test cases, some concerns have been raised about MT cost-effectiveness. Consequently, the reduction in testing costs associated with MT has become another important issue to be addressed. MR composition can address both of these problems. This technique can automatically generate new MRs by composing existing ones, thereby reducing the number of follow-up test cases. Despite this advantage, previous studies on MR composition have empirically shown that some composite MRs have lower fault detection capability than their corresponding component MRs. To investigate this issue, we performed theoretical and empirical analyses to identify what characteristics component MRs should possess so that their corresponding composite MR has at least the same fault detection capability as the component MRs do. We have also derived a convenient, but effective guideline so that the fault detection capability of MT will most likely not be reduced after composition.

**Index Terms**—Metamorphic testing, metamorphic relation, metamorphic relation composition, test oracle, fault detection capability.



## 1 INTRODUCTION

TESTING is a prominent technique for software verification [1], [2]. This technique often requires the presence of a *test oracle* (or simply an *oracle*, which refers to some mechanism for the tester to verify the correctness of the software output). However, in many situations such as testing a complex numerical algorithm, the “expected” correct software output (i.e., the oracle) is often unavailable or infeasible to determine. This problem is known as the *oracle problem*, which refers to the situation where either an oracle does not exist, or an oracle does exist but cannot be practically used, possibly due to resource constraints.

Some approaches or techniques have been proposed to alleviate the oracle problem in testing [3]. Among them, metamorphic testing (MT) has been demonstrated by various studies to be a lightweight, yet effective technique. When applying MT, the necessary properties of the software under test are firstly identified from various sources, such as the software specification. These properties are expressed in the form of relations among software inputs and outputs, formally known as *metamorphic relations (MRs)*. Since its first publications in 1998 [4], [5], MT has been repeatedly found to be effective at alleviating the oracle problem in

---

\* Corresponding author.

- K. Qiu and Z. Zheng are with the School of of Automation Science and Electrical Engineering, Beihang University, Beijing 100191, China. E-mail: {qiukun, zhengz}@buaa.edu.cn.
- T. Y. Chen is with the Department of Computer Science and Software Engineering, Swinburne University of Technology, Hawthorn VIC 3122, Australia. E-mail: tychen@swin.edu.au.
- P.-L. Poon is with the School of Engineering and Technology, Central Queensland University, Melbourne VIC 3000, Australia. E-mail:frits

composition. On the other hand, Liu et al. [31] reported that the fault detection capability may be reduced after MR composition, but without mentioning: (a) under what situations the fault detection capability will be reduced; or (b) how to avoid such a reduction happening.

In view of the above mixed findings, this paper aims to answer the question: *In what situations is testing with a composite MR more cost-effective than testing with its component MRs?* Our analysis has discovered some desirable characteristics in the component MRs, which can be easily verified by the tester, that indicate whether or not the composite MR should be used instead for testing. These characteristics are defined in terms of the bijectivity/injectivity of the input/output mappings of the component MRs. In addition to formally proving the validity of these characteristics, we have also checked their practicality through an empirical study. Furthermore, we propose a useful guideline for a tester to better decide whether or not to use MR composition. In brief, given a pair of metamorphic relations  $MR$  and  $MR'$  that fall into a special class of MRs (as defined in Definition 1 of Section 3.1), where  $MR$  is composable with  $MR'$ , they should be used to form a composite MR if both of the following conditions are met: (a) the output mapping of  $MR$  is injective; and (b) the input mapping of  $MR'$  is a bijective mapping from the source inputs of  $MR$  to the source inputs of  $MR'$ .

The rest of this paper is structured as follows. Section 2 outlines the concept of MT and gives the motivation for this study. Section 3 provides the basic concepts and terminology, which facilitate the subsequent discussion of our theoretical analysis of MR composition, which is discussed in detail in Section 4. Section 5 complements Section 4 by discussing our empirical analysis of MR composition. Based on the analyses in Sections 4 and 5, Section 6 presents our general guideline for MR composition and an analysis of its applicability. This section also discusses some related works. Finally, Section 7 summarizes and concludes the paper.

## 2 BACKGROUND

### 2.1 Metamorphic Testing (MT)

MT is a lightweight, elegant, and effective technique for alleviating the oracle problem. An intuition underlying MT is as follows: Even if we cannot verify the correctness of an *individual* output, it may still be possible to use the relations among multiple inputs and outputs for program verification.

**Example 1 (Shortest Path in an Undirected Graph):** Consider an algorithm  $f$  for computing the length of the shortest path between any two nodes ( $a$  and  $b$ ) in an undirected graph  $G$ . Let: (i)  $P$  denote an implementation of  $f$ ; (ii)  $P[G, a, b]$  denote the length of the shortest path between  $a$  and  $b$  in  $G$ , which is computed by  $P$ ; and (iii)  $f(G, a, b)$  denote the expected (and correct) length of the shortest path. When  $G$  contains many nodes and edges, for any  $a$  and  $b$  in  $G$ , it is resource-intensive and time-consuming to compute  $f(G, a, b)$  in a brute force manner for comparing with  $P[G, a, b]$ . This is because the computational complexity is of factorial order of the number of nodes in  $G$ .

With MT, this tedious verification task can be alleviated by checking, for example, two properties, which are expressed as metamorphic relations  $MR_1$  and  $MR_2$ , as follows:

- $MR_1$ : If  $a$  and  $b$  in  $G$  are swapped, then  $f(G, a, b) = f(G, b, a)$ ;
- $MR_2$ : If  $G'$  is a permutation of  $G$  with  $a'$  and  $b'$  being the permuted counterparts of  $a$  and  $b$ , respectively, then  $f(G', a', b') = f(G, a, b)$ .

With respect to  $MR_1$  and  $MR_2$ , we should have  $P[G, a, b] = P[G, b, a]$  and  $P[G, a, b] = P[G', a', b']$ . Otherwise, we can conclude that  $P$  is faulty.  $\square$

Since its first publication in 1998 [4], [5], MT has been successfully applied across a wide range of application domains and platforms. Recently, funded by the UK Engineering and Physical Sciences Research Council and the TETRACOM (TEchnology TRAnsfer in COMputing systems) EU project, a group of academics and researchers from the Department of Computing at Imperial College London (ICL) established GraphicsFuzz—a spinout company from ICL. GraphicsFuzz [16] combined fuzzing and MT to test graphics drivers. The company was acquired by Google in 2018.

A core concept of MT is the MR, which is a necessary property of a targeted function  $f$ . An MR of  $f$  is a relation over a sequence of two or more inputs  $ht_1, t_2, \dots, t_i$  and their corresponding outputs  $hf(t_1), f(t_2), \dots, f(t_i)$ , where  $n \geq 2$ . An MR can be written as  $\mathbb{R} : X \rightarrow Y$ , where  $X \times Y$  are Cartesian products of the  $n$  inputs and their corresponding  $n$  outputs. Generally, an MR can be represented as:

$$\mathbb{R}(t_1, t_2, \dots, t_i, f(t_1), f(t_2), \dots, f(t_i)).$$

Consider, for instance,  $MR_1$  in Example 1, which can be rewritten as:

$$\mathbb{R}((G, a, b), (G, b, a), f(G, a, b), f(G, b, a)).$$

Give any MR, there exists a  $k$ , where  $1 \leq k < n$ , such that:

- $t_1, t_2, \dots, t_i$  denote the *source inputs*;
- $f(t_1), f(t_2), \dots, f(t_i)$  denote the *source outputs*;
- $t_{+1}, t_{+2}, \dots, t_i$  denote the *follow-up inputs*;
- $f(t_{+1}), f(t_{+2}), \dots, f(t_i)$  denote the *follow-up outputs*.

Let  $P$  be an implementation of  $f$ . With respect to an MR, applying MT typically proceeds as follows:

- (1) Replace  $f$  with  $P$  in  $\mathbb{R}$ .
- (2) Execute  $P$  on a sequence of source inputs  $ht_1, t_2, \dots, t_i$  to obtain the corresponding sequence of source outputs  $hP[t_1], P[t_2], \dots, P[t_i]$ .
- (3) Generate a sequence of follow-up inputs  $ht_{+1}, t_{+2}, \dots, t_i$  in accordance with  $\mathbb{R}$ .
- (4) Execute  $P$  on the sequence of follow-up inputs to obtain the corresponding sequence of outputs  $hP[t_{+1}], P[t_{+2}], \dots, P[t_i]$ .
- (5) Compare the two sets of execution results with reference to  $\mathbb{R}$ . If  $\mathbb{R}$  is violated, then  $P$  is faulty.

The above steps are repeated for every identified MR.

## 2.2 MR Composition

The composition technique was originally proposed as a method of generating new MRs from existing ones [30], [31]. Example 2 explains the basic concept of MR composition.

**Example 2 (MR Composition):** Consider the following two MRs, corresponding to two well-known properties of the *sine* function:

- $MR_1$ : If  $x' = x$ , then  $\sin(x') = \sin(x)$ ;
- $MR_2$ : If  $x' = x + 2\pi$ , then  $\sin(x') = \sin(x)$ .

We can compose  $MR_1$  and  $MR_2$  together to form a composite metamorphic relation  $MR_{12}$ , that is  $MR_1(MR_2)$ .  $MR_{12}$  is formally expressed as: If  $x' = (x + 2\pi)$ , then  $\sin(x') = \sin(x)$ .  $\square$

Example 2 above shows that MR composition can generate new MRs from existing ones. If we only generate test cases from  $MR_{12}$  (and ignore  $MR_1$  and  $MR_2$ ) for testing, the testing cost is obviously reduced. In Example 2, testing with both  $MR_1$  and  $MR_2$  involves three program executions: one for  $\sin(x)$ , one for  $\sin(x + 2\pi)$ , and one for  $\sin(x)$ . On the other hand, testing with  $MR_{12}$  only involves two executions: one for  $\sin(x)$  and another for  $\sin(x + 2\pi)$ . Thus, if we only consider the testing cost, testing with  $MR_{12}$  alone is definitely preferable to testing with both  $MR_1$  and  $MR_2$ . However, beyond savings in testing costs, we should also compare the fault detection capability of  $MR_{12}$  with that of  $MR_1$  and  $MR_2$ . This leads to the research question of this paper: **(RQ) Will testing a composite MR (e.g.,  $MR_{12}$ ) have the same chance of detecting faults when compared with testing its component MRs (e.g.,  $MR_1$  and  $MR_2$ )?**

Previous studies on MR composition do not provide a definite answer to RQ. For instance, the case study reported by Dong et al. [30] has provided a "Yes" answer to RQ. On the other hand, the study by Liu et al. [31] reported that this is not necessarily the case. To date, to the best of our knowledge, no systematic studies have been conducted to address our RQ. In view of this, we perform a theoretical analysis with the intention of providing a definite answer.

## 3 IMPORTANT CONCEPTS AND TERMINOLOGY

Before we present our theoretical analysis, we first formalize some important definitions and concepts.

### 3.1 Metamorphic Relations (MRs)

In this paper, without loss of generality, we assume that the targeted function (or algorithm) involves one single input and one single output. However, generalizing our results to functions with multiple inputs and outputs is straightforward.

Because composing any two MRs into their corresponding composite MR is not always feasible, our work only considers the special class of MRs defined in this subsection. Before formally presenting this special class of MRs, let us revisit some basic concepts of mapping.

### Basic Concepts of a Function

Let  $f: A \rightarrow B$  be a function (or mapping) from  $A$  to  $B$ . Here:

- $A$  is referred to as the *domain* of  $f$ ;
- $B$  is referred to as the *codomain* of  $f$ ;
- If  $f(a) = b$ , then  $b$  is referred to as the *image* of  $a$ , and  $a$  is referred to as the *preimage* of  $b$ ;
- $f$  is said to be *injective* if  $\forall a, a' \in A, f(a) = f(a')$  implies  $a = a'$ ;
- $f$  is said to be *surjective* if  $\forall b \in B, \exists a \in A$  such that  $f(a) = b$ ;
- $f$  is said to be *bijective* if  $f$  is both injective and surjective;
- For any  $S \subseteq A$ ,  $f(S)$  denotes the set of the images of elements in  $S$  under the function  $f$ , such that:

$$f(S) = \bigcup_{s \in S} f(s);$$

- $f(A)$  is referred to as the *range* of  $f$ , and  $f(A) \subseteq B$ .

The specific class of MRs considered in our study is defined as follows:

### Definition 1. A Special Class of Metamorphic Relations (MRs)

Let

- $f: T \rightarrow R$  be a targeted function;
- $I: T \rightarrow T'$  (where  $T \subseteq T'$ , and  $T' = I(T) \cup T$ ) be a mapping that takes in a source input and generates a follow-up input for  $f$ ;
- $O: R \rightarrow R'$  (where  $R = f(T)$  and  $R' = O(R) \cup R$ ) be a mapping that takes in a source output (i.e.,  $f(t)$ ) and generates a follow-up output.

A metamorphic relation  $MR$  is a necessary property of  $f$ .  $MR$  is formally expressed as follows:

$$MR: \text{Formally Expressed as } \forall t \in T, \exists t' \in T' \text{ s.t. } t' = I(t) \text{ and } \forall r \in R, \exists r' \in R' \text{ s.t. } r' = O(r) \text{ and } f(t) = r \text{ and } f(t') = r'.$$

$t$  denote an input of  $f$ . Then, for any  $t \in T$ ,  $t$  is of the form of  $hG, a, bi$ . Additionally,

- $MR_1: \exists t \in T_1 (f(I_1(t)) = O_1(f(t)))$ , where  $T_1 = T$ ,  $I_1(hG, a, bi) = hG, b, ai$ , and  $O_1(x) = x$ .
- $MR_2: \exists t \in T_2 (f(I_2(t)) = O_2(f(t)))$ , where  $T_2 = T$ ,  $I_2(\cdot)$  is a permutation function that takes in  $hG, a, bi$  and permutes  $G, a$ , and  $b$  according to a certain pattern, such that  $I_2(hG, a, bi) = hG', a', b'i$ , where  $G'$  is the permuted  $G$ ,  $a'$  and  $b'$  are the counterparts of  $a$  and  $b$ , respectively, and  $O_2(x) = x$ .

Next, consider Example 2. Let  $T$  and  $R$  be the domain and codomain of the *sine* function, respectively; we write *sine* as  $f$  and  $t$  as the input of  $f$ . Then:

- $MR_3: \exists t \in T_3 (f(I_3(t)) = O_3(f(t)))$ , where  $T_3 = T$ ,  $I_3(t) = t$ , and  $O_3(x) = x$ ;
- $MR_4: \exists t \in T_4 (f(I_4(t)) = O_4(f(t)))$ , where  $T_4 = T$ ,  $I_4(t) = t + 2\pi$ , and  $O_4(x) = x$ .

### 3.2 Composable MR and Composite MR

Below we first define a composable MR, which will facilitate the definition of a composite MR.

#### Definition 2. Composable MR

Let

- $f: T \rightarrow R$  be a targeted function;
- $MR$  and  $MR'$  be two MRs of  $f$ .

$MR$  is said to be *composable* with  $MR'$  if:

- $I(MR) \subseteq T$ , that is, the range of  $I$  is a subset of the domain of  $I'$ ;
- $O(MR) \subseteq R$  (where  $R = f(T)$  and  $R' = f(T')$ ), that is, the range of  $O$  is a subset of the domain of  $O'$ .

For the rest of the paper, we use subscripts to link an MR and its related components. For instance, in Definition 2 above,  $T, I, I',$  and  $O, O'$  denote the set of source inputs, input mapping, and output mapping corresponding to  $MR, MR'$ , respectively.

Refer to Example 1. It can be deduced that:

- $I_1(T_1) = T_1 = T_2$  because (i)  $I_1(T_1) = T_1$  and (ii)  $T_1 = T_2 = T$ ;
- $O_1(R_1) = R_1 = R_2$  because (i)  $R_1 = R_2$  as  $f(T_1) = f(T_2)$  and (ii)  $O_1(R_1) = R_1$  because  $O_1(x) = x$ ;
- $I_2(T_2) = T_2 = T_1$  because (i)  $I_2(T_2) = T_2$  as  $I_2$  is a permutation function and (ii)  $T_1 = T_2 = T$ ;
- $O_2(R_2) = R_2 = R_1$  because (i)  $R_2 = R_1$  as  $f(T_2) = f(T_1)$  and (ii)  $O_2(R_2) = R_2$  as we have  $O_2(x) = x$ .

With (a) and (b), it follows after Definition 2 that  $MR_2$  is composable with  $MR_1$ . Similarly, with (c) and (d), we conclude that  $MR_1$  is composable with  $MR_2$ . It should, however, be noted that for any two metamorphic relations  $MR$  and  $MR'$ , if  $MR$  is composable with  $MR'$ , it is not necessary that  $MR'$  is also composable with  $MR$ .

We next formally define the construction of a composite MR:

#### Definition 3. Composite MR and Component MR

Let

- $f: T \rightarrow R$  be a targeted function;
- $MR$  and  $MR'$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR'$ .

The *composite MR* (denoted by  $MR \circ MR'$ ), formed by composing  $MR$  with  $MR'$ , is a necessary property of  $f$ .  $MR \circ MR'$  is formally expressed as follows:

$$\exists t \in T (f(I \circ I'(t)) = O \circ O'(f(t))),$$

where

- $T' = T$ ;
- $I \circ I'(t) = (I \circ I')(t) = I(I'(t))$ ;
- $O \circ O'(f(t)) = (O \circ O')(f(t)) = O(O'(f(t)))$ .

We write  $MR \circ MR' = MR \circ MR'$  or  $MR(MR')$ . Also, we refer to  $MR$  and  $MR'$  as the *component metamorphic relations* of  $MR \circ MR'$ .

Refer to Example 2. It is straightforward to conclude that  $MR_1$  and  $MR_2$  are composable with each other according to Definition 2. We write  $f(t)$  as  $\sin(t)$ . Then, with respect to  $MR_1$  and  $MR_2$ , there are two possible composite MRs.

- $MR_{12} = MR_1(MR_2) = MR_1 \circ MR_2$ :

$$\exists t \in T_{12} (f(I_{12}(t)) = O_{12}(f(t))),$$

where  $T_{12} = T_2 = T$ ,  $I_{12}(t) = I_1(I_2(t)) = t + 2\pi$  and  $O_{12}(x) = O_1(O_2(x)) = x$ .

- $MR_{21} = MR_2(MR_1) = MR_2 \circ MR_1$ :

$$\exists t \in T_{21} (f(I_{21}(t)) = O_{21}(f(t))),$$

where  $T_{21} = T_1 = T$ ,  $I_{21}(t) = I_2(I_1(t)) = t + 2\pi$ , and  $O_{21}(x) = O_2(O_1(x)) = x$ .

It should be noted that, when applying MT, the tester is not required to explicitly specify the composite MR in the format above: The task of composing composable MRs can be automated through programming in accordance with Definition 3. This automation can be implemented as follows. Two test scripts can be written — one calling function  $I$  and the other calling function  $I'$ . If  $t$  is an input to  $I$ , then the returned value from  $I$  is used as an input to  $I'$ . In this way,  $t$  and the returned value from  $I$  form a pair of source and follow-up inputs for the composite  $MR \circ MR'$ . Similarly, the pair of source and follow-up outputs for  $MR \circ MR'$  could be obtained by first executing a test script to call the function  $O$ , followed by executing another test script to call the function  $O'$ . Here, the functions  $I, I', O, O'$  are implemented according to  $MR$  and  $MR'$ .

When composing more than two MRs, the resultant composite MR can also be automatically obtained by recursively applying Definition 3. Since  $I$  and  $O$  are mappings, the composition of  $I$ s and the composition of  $O$ s are associative, that is:  $I \circ (I \circ I') = (I \circ I) \circ I'$  and  $O \circ (O \circ O') = (O \circ O) \circ O'$ . Since an MR is defined in terms of its own  $I$  and  $O$ , therefore, the composition of MRs is also associative. For example,  $MR \circ (MR \circ MR') = (MR \circ MR) \circ MR'$ . The

order of composition, however, is important. For instance,  $MR \circ MR$  may not be the same as  $MR \circ MR$ . In other words, the composition of MRs is not commutative.

### 3.3 Evaluation of Fault Detection Capability

The evaluation and comparison of the fault detection capabilities of two different MRs (regardless of whether or not they are composite) requires some metric. We defined two—one qualitative, and one quantitative—to use in our study. To facilitate the definition of these two metrics, we first present the following definition:

#### Definition 4. Satisfiability of a Set of Source Inputs for an MR

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  be a metamorphic relation of  $f$  with  $T$ ,  $I$ , and  $O$  being its set of source inputs, input mapping, and output mapping, respectively;
- $S$  be a nonempty subset of  $T$ .

After executing all elements of  $S$  with  $P$ ,

- $S$  is said to *satisfy*  $MR$ , if all elements of  $S$  satisfy  $MR$ , that is,

$$\exists t \in S (O(P[t]) = P[I(t)]);$$

- $S$  is said to *violate*  $MR$ , if all elements of  $S$  violate  $MR$ , that is,

$$\exists t \in S (O(P[t]) \neq P[I(t)]).$$

With Definition 4, a qualitative metric is defined for measuring the fault detection capability of an MR as follows:

#### Definition 5. Satisfiability of an MR

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  be a metamorphic relation of  $f$  with  $T$  being its set of source inputs.

With respect to  $P$ , if  $T$  satisfies  $MR$ ,  $MR$  is said to be *satisfiable*; otherwise,  $MR$  is said to be *violative*.

Let  $T'$  denote the set of all elements in  $T$  that violate  $MR$ . In this case,  $T'$  is referred to as the *set of violative source inputs* of  $MR$ . Obviously,  $T' = \emptyset$ , iff  $MR$  is *satisfiable*.

In Definition 5 above, given an implementation  $P$  and an MR, it is equivalent to say that  $P$  violates (or satisfies) the MR, when the MR is violative (or satisfiable).

Together Definitions 4 and 5 allow us to define the following quantitative metric for measuring the fault detection capability of an MR:

#### Definition 6. Fault Detection Rate of an MR

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  be a metamorphic relation of  $f$  with  $T$  being its set of source inputs and  $T'$  being its set of violative source inputs.

Let  $\theta$  denote the *fault detection rate* of  $MR$  with respect to  $P$ . Then,  $\theta$  is defined as follows:

$$\theta = \frac{|T'|}{|T|},$$

where  $|T'|$  and  $|T|$  denote the sizes of  $T'$  and  $T$ , respectively.

The satisfiability of an MR (Definition 5) indicates whether or not a program under test can be revealed as faulty by this MR. Furthermore, the fault detection rate (Definition 6) indicates *how likely* it is that an MR will reveal a fault in  $P$  with only one *random* source input. Larger fault detection rates indicate higher fault detection capabilities.

## 4 THEORETICAL ANALYSIS OF FAULT DETECTION CAPABILITY

Given an implementation  $P$ ,  $MR_1$ , and  $MR_2$ , there are four possible scenarios:

- (1) Both  $MR_1$  and  $MR_2$  are satisfiable;
- (2)  $MR_1$  is satisfiable and  $MR_2$  is violative;
- (3)  $MR_1$  is violative and  $MR_2$  is satisfiable;
- (4) Both  $MR_1$  and  $MR_2$  are violative.

For each of the above scenarios, we analyze the fault detection capability of  $MR_1$ .

### 4.1 Scenario 1

In this scenario, both  $MR_1$  and  $MR_2$  are satisfiable, that is,  $\theta_1 = \theta_2 = 0$ . Although we intuitively expect  $\theta$  to be 0, let us formally prove it (Theorem 1). This proof needs the following lemma.

#### Lemma 1.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR_1$  and  $MR_2$  be two MRs of  $f$ ;
- $MR_1$  be composable with  $MR_2$ ;
- $S$  be a nonempty subset of  $T$ .

If  $S$  satisfies  $MR_1$  and  $I(S)$  satisfies  $MR_2$ , then  $S$  satisfies  $MR_2$ .

*Proof (Lemma 1).* Assume that  $S$  satisfies  $MR_1$  and  $I(S)$  satisfies  $MR_2$ . It follows after Definition 4 that

$$\exists t \in S (O(P[t]) = P[I(t)]), \quad (1)$$

and

$$\exists t' \in I(S) (O(P[t']) = P[I(t')]). \quad (2)$$

By the definition of  $I(S)$ , for any  $t' \in I(S)$ , there exists a  $t \in S$  such that  $t' = I(t)$ ; and for any  $t \in S$ , there exists a  $t' \in I(S)$  such that  $t' = I(t)$ . Therefore, Eq. 2 can be rewritten as follows:

$$\exists t \in S (O(P[I(t)]) = P[I(I(t))]). \quad (3)$$

Immediately after Eqs. 1 and 3, we have

$$\exists t \in S (O(O(P[t])) = P[I(I(t))]). \quad (4)$$

It follows after Definition 4 that  $S$  satisfies  $MR$ .  $\square$

Now we are ready to present Theorem 1 and its proof.

### Theorem 1.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  and  $MR$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR$ .

If  $MR$  and  $MR$  are satisfiable, then their composite metamorphic relation,  $MR$ , is also satisfiable.

*Proof (Theorem 1).* Assume that  $MR$  and  $MR$  are both satisfiable. From Definition 5, we immediately have  $T$  satisfies  $MR$  and  $T$  satisfies  $MR$ . Since  $MR$  is composable with  $MR$ , it follows after Definition 2 that  $I(T) = T$ . Since  $T$  satisfies  $MR$ ,  $I(T)$  also satisfies  $MR$ .

It follows from Lemma 1 that, since  $T$  satisfies  $MR$  and  $I(T)$  satisfies  $MR$ ,  $T$  satisfies  $MR$ . Because  $T = T$  (Definition 3), therefore, it follows after Definition 5 that  $MR$  is satisfiable, that is,  $\theta = 0$ .  $\square$

**Implication.** Theorem 1 states that, if an implementation  $P$  under test does not violate any two component MRs ( $MR$  and  $MR$ ),  $P$  will also not violate any composite MR constructed from  $MR$  and  $MR$ .

## 4.2 Scenario 2

In this scenario,  $MR$  is satisfiable ( $\theta = 0$ ) and  $MR$  is violative ( $\theta > 0$ ). Before analyzing the fault detection capability of  $MR$ , we first introduce the following lemma to facilitate the proof of Theorem 2.

### Lemma 2.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  and  $MR$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR$ ;
- $S$  be a nonempty subset of  $T$ .

Suppose that  $S$  violates  $MR$  and  $I(S)$  satisfies

$MR$ . If  $O$  is an injective mapping, then  $S$  violates  $MR$ .

*Proof (Lemma 2).* Since  $S$  violates  $MR$ , it follows after Definition 4 that

$$\exists t \in S (O(P[t]) \notin P[I(t)]). \quad (5)$$

Because  $I(S)$  satisfies  $MR$ , it follows after Definition 4 that

$$\exists t' \in I(S) (O(P[t']) = P[I(t')]). \quad (6)$$

By the definition of  $I(S)$ , for any  $t' \in I(S)$ , there exists a  $t \in S$  such that  $t' = I(t)$ ; and for any  $t \in S$ , there exists a  $t' \in I(S)$  such that  $t' = I(t)$ . Therefore, Eq. 6 can be rewritten as follows:

$$\exists t \in S (O(P[I(t)]) = P[I(I(t))]). \quad (7)$$

Assume that  $O$  is an injective mapping. Immediately after Eqs. 5 and 7, we have

$$\exists t \in S (O(O(P[t])) \notin P[I(I(t))]). \quad (8)$$

Therefore,  $S$  violates  $MR$ .  $\square$

With Lemma 2, we can now introduce Theorem 2 and its proof.

### Theorem 2.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  and  $MR$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR$ .

Suppose that  $MR$  is satisfiable ( $\theta = 0$ ) and  $MR$  is violative ( $\theta > 0$ ). If  $O$  is an injective mapping, then  $MR$  is violative and  $\theta = \theta$ .

*Proof (Theorem 2).* To determine  $\theta$ , we need to know the set of source inputs ( $T$ ) and the set of violative source inputs ( $T$ ) of  $MR$ . It follows from Definition 3 that  $T = T$ . In what follows, we will prove that, if  $O$  is an injective mapping, then  $T = T$ .

Since  $MR$  is violative, we have  $T = T \setminus \overline{T}^{-1}$ , where  $T \notin$ . Since  $MR$  is composable with  $MR$ , we have  $I(T) = T$ . In turn, we have  $I(T) = T$  and  $I(\overline{T}) = T$ . Since  $MR$  is satisfiable, it follows after Definition 5 that  $T$  satisfies  $MR$ . Therefore, we have

$$I(T) \text{ satisfies } MR; \quad (i)$$

and

$$I(\overline{T}) \text{ satisfies } MR, \text{ if } I(\overline{T}) \notin; \quad (ii)$$

Next, let us assume that  $O$  is injective. Since we have (i) above and  $T$  violates  $MR$ , it follows from Lemma 2 that

1. In this paper, we use  $\overline{T^v}$  to denote the complementary set of  $T^v$  over  $T$ . For instance,  $\overline{T^y}$  is the complementary set of  $T^y$  over  $T_y$ .

$$T \text{ violates } MR. \quad (iii)$$

Furthermore, we have

$$T = T, \quad (iv)$$

because  $T$  contains all the elements that violate  $MR$ .

Next, we consider the following two exhaustive cases:

Case (a):  $\overline{T} \neq \emptyset$ ;

Since we have (ii) and  $\overline{T}$  satisfies  $MR$ , it follows from Lemma 1 that  $\overline{T}$  satisfies  $MR$ . Furthermore, since we have (iii) and  $T = T = T \setminus [\overline{T}]$ , therefore, we have  $T = T$ .

Case (b):  $\overline{T} = \emptyset$ ;

Since we have (iv),  $T = T$  (where  $T = T$ ), and  $T = T \setminus [\overline{T} = T]$ , therefore, we have  $T = T$ .

In view of the above two cases, regardless of whether or not  $\overline{T}$  is empty, we have  $T = T$ . Then, it follows after Definition 6 that,

$$\theta = \frac{|\overline{T}|}{|T|} = \frac{|\overline{T}|}{|T|} = \theta > 0.$$

In other words,  $MR$  is violative.  $\square$

**Implication:** Theorem 2 gives a sufficient condition for  $MR$  and  $MR$  having the same fault detection capability if  $MR$  is satisfiable.

A previous study [31] reported that composing some "loose" MRs may result in a composite MR with a lower fault detection capability. However, that study has not formally defined the meaning of "loose" MRs. By means of our theoretical analysis, we found that a "loose" MR in fact refers to one whose output mapping is not injective.

### 4.3 Scenario 3

In this scenario,  $MR$  is violative ( $\theta > 0$ ) and  $MR$  is satisfiable ( $\theta = 0$ ). Before introducing Theorems 3 and 4, we need the following lemma to facilitate their proofs.

#### Lemma 3.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  and  $MR$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR$ ;
- $S$  be a nonempty subset of  $T$ .

If  $S$  satisfies  $MR$  and  $I(S)$  violates  $MR$ , then  $S$  violates  $MR$ .

*Proof (Lemma 3).* Assume that there exists a nonempty  $S \subseteq T$  such that  $S$  satisfies  $MR$  and  $I(S)$  violates  $MR$ . Since  $S$  satisfies  $MR$ , it follows after Definition 4 that

$$\exists t \in S \ (O(P[t]) \neq P[I(t))). \quad (9)$$

Because  $I(S)$  violates  $MR$ , it follows after Definition 4 that

$$\exists t' \in I(S) \ (O(P[t']) \neq P[I(t'))). \quad (10)$$

By the definition of  $I(S)$ , for any  $t' \in I(S)$ , there exists a  $t \in S$  such that  $t' = I(t)$ ; and for any  $t \in S$ , there exists a  $t' \in I(S)$  such that  $t' = I(t)$ . Therefore, Eq. 10 can be rewritten as follows:

$$\exists t \in S \ (O(P[I(t)]) \neq P[I(I(t))]). \quad (11)$$

Immediately after Eqs. 9 and 11, we have

$$\exists t \in S \ (O(O(P[t])) \neq P[I(I(t))]). \quad (12)$$

Therefore,  $S$  violates  $MR$ .  $\square$

#### Theorem 3.

Let

- $f$  be a targeted function;
- $P$  be an implementation of  $f$ ;
- $MR$  and  $MR$  be two MRs of  $f$ ;
- $MR$  be composable with  $MR$ .

Suppose that  $MR$  is violative ( $\theta > 0$ ) and  $MR$  is satisfiable ( $\theta = 0$ ). If  $I(T) = T$ , then  $MR$  is violative with  $\theta = \frac{|\overline{T}|}{|T|} > 0$  (where  $T \neq \emptyset$ ,  $T = T$ , and  $I(T) = T$ ).

*Proof (Theorem 3).* To determine  $\theta$ , we need to know the set of source inputs ( $T$ ) and the set of violative source inputs ( $\overline{T}$ ) of  $MR$ . It follows from Definition 3 that  $T = T$ . In what follows, we will prove that, if  $I(T) = T$ , then  $T = T$  (where  $T \neq \emptyset$ ,  $T = T$ , and  $I(T) = T$ ).

Since  $MR$  is violative, we have  $T = T \setminus [\overline{T}]$ , where  $T \neq \emptyset$ . Let us assume that: (a)  $I(T) = T$ ; and (b)  $T, T = T$ , such that  $I(T) = T$  and  $I(T) = \overline{T}$ . Since  $T \neq \emptyset$ , we have  $T \neq \emptyset$ ;  $\square$ , *ad*





TABLE 1: Fault Detection Rates ( $\theta$ ) of  $MR$  in Different Testing Scenarios

Scenario	$\theta_x$	$\theta_y$	$\max\{\theta_x, \theta_y\}$	$\theta_{xy}$
1	= 0	= 0	= 0	$\theta_{xy} = 0$ (Theorem 1)
2	= 0	> 0	= $\theta_y$	If $O_x$ is injective, then $\theta_{xy} = \theta_y$ (Theorem 2)
3	> 0	= 0	= $\theta_x$	If $I_y(T_y) = T_x$ , then $\theta_{xy} > 0$ (Theorem 3)
				If $I_y(T_y) = T_x$ and $I_y$ is bijective, then $\theta_{xy} = \theta_x$ (Theorem 4)
4	> 0	> 0	Varies in different situations	If $I_y(T_y) = T_x$ and $O_x$ is injective, then it is very likely to have $\theta_{xy} > 0$

$$A_3 \cap T = \emptyset. \quad (\text{viii})$$

Let us assume that  $MR$  is satisfiable, that is,  $T = \emptyset$  or  $\theta = 0$ . Immediately, with (vii) and (viii) above, we have  $A_2 = A_3 = \emptyset$ . Because  $T = A_1 \cap A_2 = A_1 \cap \emptyset = \emptyset$ , we have  $T = \emptyset$ . Since  $T = A_3 \cap A_4 = \emptyset \cap A_4 = \emptyset$ , we have  $T = \emptyset$ . Furthermore, since  $T = T \cap T = \emptyset \cap \emptyset = \emptyset$ , we have  $T = \emptyset$  and  $T = \emptyset$ . Since  $I(T) = \emptyset$  and  $I(T) = \emptyset$  by definition, we have

$$I(T) = \emptyset, \quad (\text{ix})$$

and

$$I(\emptyset) = \emptyset. \quad (\text{x})$$

Since (ix) and (x) follow after the assumption that  $\theta = 0$ , they are necessary relations for  $\theta$  to be 0. Relation (ix) implies that, for every violating source test input for  $MR$ , its corresponding follow-up input must be a violating source input for  $MR$ . Similarly, relation (x) implies that, for every non-violating source input for  $MR$ , its corresponding follow-up input must be a non-violating source input for  $MR$ . Obviously, these two relations are very tight and restrictive and, hence, they are unlikely to be satisfied simultaneously. Since the situation of  $\theta = 0$  requires the simultaneous satisfaction of relations (ix) and (x) (which is very rare, as explained above), it can be comfortably concluded that the situation of  $\theta > 0$  is very likely to occur. In summary, the above theoretical analysis has showed that our hypothesis will be strongly held because of two very tight and restrictive relations (ix) and (x). We performed an empirical study to support the above theoretical analysis for Scenario 4. Details will be given in Section 5.

Table 1 summarizes our theoretical analysis on the fault detection rates of  $MR$  in the four different testing scenarios.

## 5 EMPIRICAL ANALYSIS OF FAULT DETECTION CAPABILITY

In the first three testing scenarios (Scenarios 1, 2, and 3) discussed in Sections 4.1, 4.2, and 4.3, we are able to obtain a definite answer after a theoretical analysis. In other words, for each of those three scenarios, we have found the characteristics that component MRs should possess to guarantee that a composite MR has the same chance of detecting the faults as its component MRs do. On the other hand, Scenario 4 is too complicated to have a definite answer solely based on a theoretical analysis. Nevertheless, our theoretical analysis of Scenario 4 has led to a hypothesis, as stated at

the beginning of Section 4.4, which was further verified by an empirical study to be discussed in this section. We next discuss the settings and observations of our empirical study.

### 5.1 Subject Programs

Our empirical study involved the following four subject programs:

- *TriangleSquare (TSQ)*. It accepts three numbers, corresponding to the three edges of a triangle, and calculates its area if a legitimate triangle can be formed [30].
- *SparseMatrixMultiplication (SMM)*. It accepts two sparse matrices<sup>2</sup> as inputs and computes their product matrix [30].
- *Dnapars (DNA)*. It is commonly used in bioinformatics [8]. It takes in a matrix containing a set of species' DNAs and generates an evolution tree.
- *K-Nearest Neighbors (KNN)*. It is a machine learning classifier algorithm, which takes in a training data set and a testing data, and then predicts the label for the latter based on the former [19].

Table 2 gives more details about these subject programs, in terms of their inputs, outputs, and the approaches we used to assert two equivalent outputs.

### 5.2 Experimental Procedures

We applied the following steps to each subject program:

- (1) For each identified MR, manually check its compliance with Definition 1. This check is required because Definition 1 specifies a special (but common) class of MRs that our study has assumed.
- (2) For any tuple of two MRs ( $MR_1, MR_2$ ), use Definition 2 to manually check whether  $MR_1$  is composable with  $MR_2$ . If yes, then check whether: (a)  $O_1$  is injective; and (b)  $I_2(T_1) = T_2$ . If yes to both (a) and (b), then generate the composite MR ( $MR_1 \circ MR_2$ ) from  $MR_1$  and  $MR_2$ .
- (3) Apply mutation analysis and random testing to estimate individual fault detection rates for  $MR_1, MR_2$ , and  $MR_1 \circ MR_2$ .

### 5.3 Component and Composite MRs

Table 3 lists all the MRs used in our study; they were sourced from previous MT-related studies [8], [19], [30]. All these MRs were then checked and confirmed to comply with Definition 1 (see step 1). There were 23 ( $7 + 7 + 6 + 3$ ) such MRs for the four subject programs. The 3rd column of this

2. A sparse matrix is a matrix in which most of the elements are zero.

TABLE 2: Inputs, Outputs, and Equivalent Output Assertions of Subject Programs

Subject Programs	Inputs ( )	Outputs ( ( ))	Assertion of Equivalent Outputs
<i>TSQ</i>	$\text{Inputs} = \langle a, b, c \rangle$ , where $a, b$ , and $c$ denote the three edges of a triangle	If the three edges form a legitimate triangle, then $\text{Output} = \text{Area}$ , where $\text{Area}$ denotes the area of the triangle	Let $\text{Area}$ and $\text{Area}'$ denote two outputs. Then, $\text{Area} = \text{Area}'$ iff $ \text{Area} - \text{Area}'  < \epsilon^\dagger$ .
<i>SMM</i>	$\text{Inputs} = \langle A, B \rangle$ , where $A$ and $B$ denote two sparse matrices in the Compressed Sparse Row (CSR) format	$\text{Output} = C$ , where $C$ denotes the product of $A$ and $B$	Let $C = [c_{ij}]$ and $C' = [c'_{ij}]$ denote two outputs. Then, $C = C'$ iff for all $i, j$ , $ c_{ij} - c'_{ij}  < \epsilon^\dagger$ .
<i>DNA</i>	$\text{Inputs} = X$ , where $X$ denotes the species' DNAs and is in the format of a $(N \times L)$ matrix ( $N$ = number of species; $L$ = length of a DNA sequence)	$\text{Output} = \langle \text{Length}, \text{Tree} \rangle$ , where $\text{Length}$ and $\text{Tree}$ denote the length and the structural description of the generated evolution tree, respectively	Let $\langle \text{Length}, \text{Tree} \rangle$ and $\langle \text{Length}', \text{Tree}' \rangle$ denote two outputs. Then, $\langle \text{Length}, \text{Tree} \rangle = \langle \text{Length}', \text{Tree}' \rangle$ iff $ \text{Length} - \text{Length}'  < \epsilon^\dagger$ and $\text{Tree}$ is identical to $\text{Tree}'$ .
<i>KNN</i>	$\text{Inputs} = \langle X, C, S \rangle$ , where $X$ denotes the attributes of the training data set in the format of a $(N \times M)$ matrix ( $N$ = number of entries of the training data set; $M$ = number of attributes of each entry), $C$ denotes the class labels of the training data with a size of $N$ , and $S$ denotes the testing item's attributes with a size of $M$ elements	$\text{Output} = c_v$ , where $c_v$ denotes the calculated class label for $S$	Let $c_v$ and $c'_v$ denote two class labels. Then, $c_v = c'_v$ iff $c_v$ and $c'_v$ are exactly the same.

( $\dagger$ )  $\epsilon$  is an extremely small value and is set to  $10^{-6}$ .

Table 2 gives the details on these MRs. For each such MR, we also explicitly list its set of source inputs (4th column), input mapping (5th column), and output mapping (6th column) in the table.

For any tuple of two MRs ( $MR_1, MR_2$ ), checks were performed to ensure that conditions (a) and (b) in step 2 were fulfilled. After checking for all the tuples of two MRs, a total of 106 ( $TSQ=42, SMM=42, DNA=16$ , and  $KNN=6$ ) eligible pairs of component MRs were found, thereby resulting in the construction of 106 composite MRs. As an example, let us consider  $MR_1$  and  $MR_2$  of *TSQ* in Table 3. We write the function corresponding to *TSQ* as  $f$ . It can be deduced that: (a)  $T_2 = I_1(T_1) = T_1 = I_2(T_2) = T$  because both  $I_1$  and  $I_2$  are bijective mappings from  $T$  to  $T$ ; and (b)  $R_2 = O_1(R_1) = R_1 = O_2(R_2) = f(T)$  because both  $O_1$  and  $O_2$  are bijective mappings from  $f(T)$  to  $f(T)$ . It then follows from Definition 2 that  $MR_1$  and  $MR_2$  are composable with each other. Therefore,  $MR_{12}$  and  $MR_{21}$  were formed and tested in our study.

#### 5.4 Measurement of Fault Detection Rates and Generation of Mutants

According to Definition 6, the fault detection rate of an MR with respect to a program  $P$  is the ratio of  $\sum_j V_j$  (the size of the MR's set of violative source inputs) to  $\sum_j S_j$  (the size of the MR's set of source inputs). However, because the size of  $T$  is often very large, therefore it is practically infeasible to conduct exhaustive testing to determine the value of  $T$ . In turn, we cannot compute an MR's fault detection rate based on Definition 6. Therefore, in this study, we used the following equation as the "estimator" of an MR's fault detection rate ( $\theta$ ):

$$\theta = \frac{N_v}{N_s}$$

where  $N_v$  denotes the number of tested source inputs that caused violations to a given MR, and  $N_s$  denotes the total number of source inputs used in testing.

Mutation analysis [32] has long been used in MT to evaluate the fault detection rate of MRs (e.g., in [19]). Thus, we also used the mutation technique (together with random testing) to estimate the fault detection rates of component and composite MRs. Table 4 shows the mutants of the four subject programs with injected faults. We randomly generated 10 000 ( $N = 10\,000$ ) source inputs for each of the four subject programs.

For each mutant of every subject program, we then performed two operations: (a) used each identified component MR and each constructed composite MR to generate a separate set of follow-up inputs from the set of source inputs (with a size of 10 000); and (b) executed these source inputs and follow-up inputs with the mutants and checked for violations to MRs.

#### 5.5 Experimental Observations

Table 5 shows the estimated fault detection rates of component and composite MRs when  $\theta > 0$  and  $\theta > 0$  (a total of 108 such cases). The complete set of experimental data of the estimated fault detection rates for the four subject programs are given in Tables 7, 8, and 9 in the Appendix (as a separate file).

For ease3

Figure 4 shows the estimated fault detection rates of component and composite MRs when  $\theta > 0$  and  $\theta > 0$  (a total of 108 such cases). The complete set of experimental data of the estimated fault detection rates for the four subject programs are given in Tables 7, 8, and 9 in the Appendix (as a separate file).

TABLE 3: List of MRs Used in the Empirical Study

Sub. Pro.	MRs	Description	$T$	Input Mapping	Output Mapping
TSQ	$MR_1$	If $\langle a', b', c' \rangle = \langle b, a, c \rangle$ , then $' =$	$T_1 = \mathcal{T}$	$I_1(\langle a, b, c \rangle) = \langle b, a, c \rangle$	$O_1( ) =$
	$MR_2$	If $\langle a', b', c' \rangle = \langle a, c, b \rangle$ , then $' =$	$T_2 = \mathcal{T}$	$I_2(\langle a, b, c \rangle) = \langle a, c, b \rangle$	$O_2( ) =$
	$MR_3$	If $\langle a', b', c' \rangle = \langle c, b, a \rangle$ , then $' =$	$T_3 = \mathcal{T}$	$I_3(\langle a, b, c \rangle) = \langle c, b, a \rangle$	$O_3( ) =$
	$MR_4$	If $\langle a', b', c' \rangle = \langle 2a, 2b, 2c \rangle$ , then $' = 4$	$T_4 = \mathcal{T}$	$I_4(\langle a, b, c \rangle) = \langle 2a, 2b, 2c \rangle$	$O_4( ) = 4$
	$MR$	If $\langle a', b', c' \rangle = \langle \sqrt{2b^2 + 2c^2 - a^2}, b, c \rangle$ , then $' =$	$T = \mathcal{T}$	$I(\langle a, b, c \rangle) = \langle \sqrt{2b^2 + 2c^2 - a^2}, b, c \rangle$	$O( ) =$
	$MR$	If $\langle a', b', c' \rangle = \langle a, \sqrt{2a^2 + 2c^2 - b^2}, c \rangle$ , then $' =$	$T = \mathcal{T}$	$I(\langle a, b, c \rangle) = \langle a, \sqrt{2a^2 + 2c^2 - b^2}, c \rangle$	$O( ) =$
	$MR$	If $\langle a', b', c' \rangle = \langle a, b, \sqrt{2a^2 + 2b^2 - c^2} \rangle$ , then $' =$	$T = \mathcal{T}$	$I(\langle a, b, c \rangle) = \langle a, b, \sqrt{2a^2 + 2b^2 - c^2} \rangle$	$O( ) =$
$MR_1$	If $\langle A'T$				

SMM

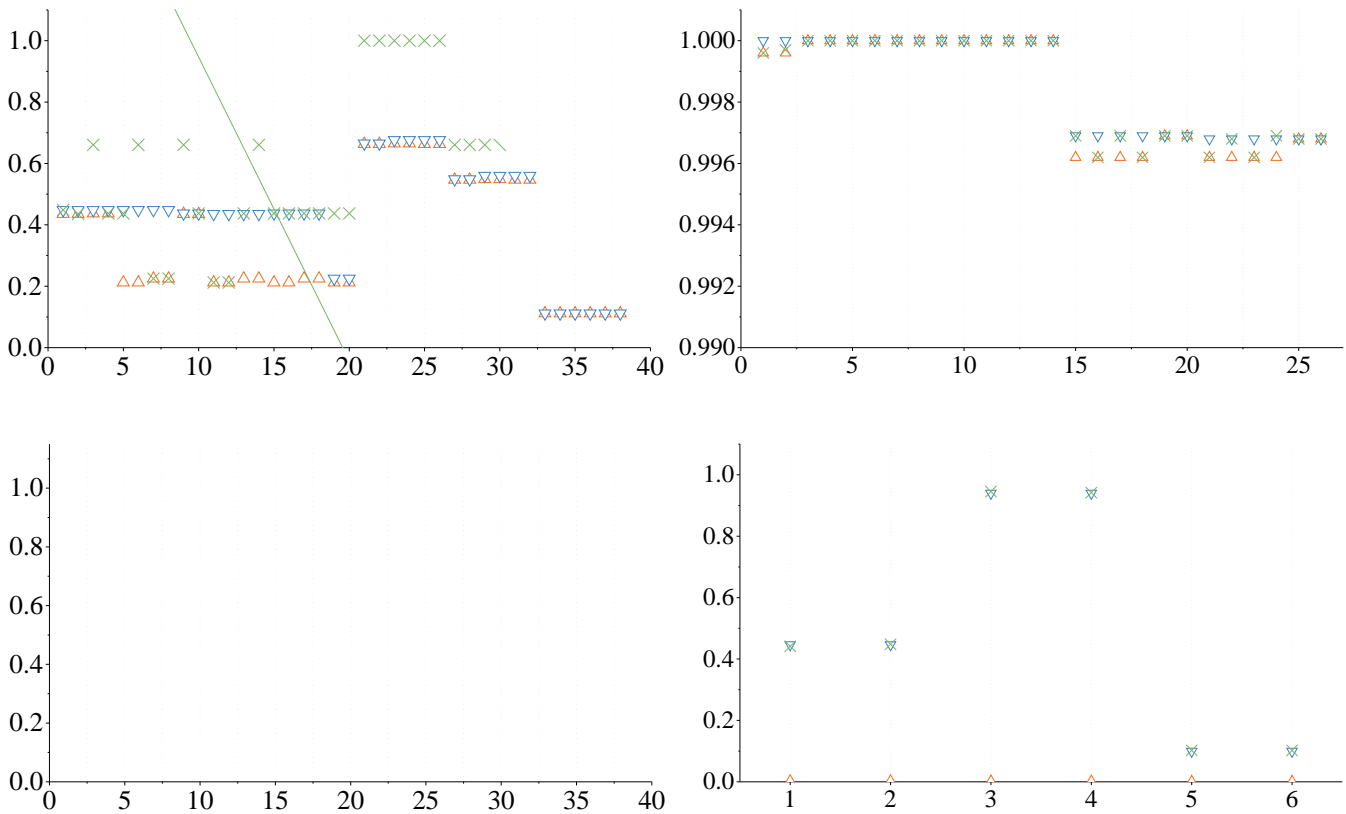


Fig. 1: Comparison among  $\min f\theta$ ,  $\theta g$ ,  $\max f\theta$ ,  $\theta g$ , and  $\theta$ .

step (a), the composition of MRs should not be complex. However, as a precaution to ensure that the generated composite MRs were valid with respect to their corresponding subject programs, we performed two tasks. Firstly, we conducted a desk check on these composite MRs, during which we detected no abnormality. Secondly, we tested all the four subject programs (their "original" versions; not their mutants) against these composite MRs. The testing results did not reveal any MR violation. To a large extent, these two tasks provide assurance that the composite MRs were correctly generated.

**Subject program selection:** Undoubtedly, it would be desirable to have a large set of programs for our empirical study. However, using a large set of programs was prohibited due to resource constraints. Nevertheless, using these four subject programs still provided a good insight into the validity of the hypothesis as stated at the beginning of Section 4.4 because these programs cover: (a) different application domains, including numerical calculations (*TSQ* and *SMM*), bioinformatics (*DNA*), and machine learning classifiers (*KNN*); and (b) different levels of complexity (*TSQ* and *KNN* are relatively less complex in logic, whereas *SMM* and *DNA* are relatively more complex).

**Mutant generation and selection:** The mutants used for *TSQ*, *SMM*, and *DNA* were sourced from other previous studies [8], [30], whereas the mutants used for *KNN* were

generated by us (because we could not find mutants for this program from the published work). To a large extent, selecting extent,

TABLE 4: Mutants of Subject Programs

Sub. Pro.	Mu. ID	Code Change for Mutant Generation
TSQ	$\mu_1$	Swap lines 6 and 8
	$\mu_2$	Replace "p=(a+b+c)/2" by "p=(a+b+c)*2" in line 22
	$\mu_3$	Replace "/2" by "*2" in lines 32, 40, and 48
	$\mu_4$	Replace "(math.sqrt(3.0)*a*a)/4.0" by "(math.sqrt(3.0)*a*a)/2.0" in line 102
SMM	$\mu_1$	Replace "n" by "1" in line 43
	$\mu_2$	Replace "c[nz] = aij * b[k]" by "c[nz] = aij" in line 52
	$\mu_3$	Replace "c[nz] = aij * b[k]" by "c[nz] = b[k]" in line 52
	$\mu_4$	Replace "c[icol] += aij*b[k]" by "c[icol] += aij" in line 56
	$\mu$	Replace "c[icol] += aij*b[k]" by "c[icol] += b[k]" in line 56
DNA	$\mu_1$	Replace "ns=1<<G" by "ns=1<<C" in line 720 in file "seq.c"
	$\mu_2$	Replace "ally[alias[i-1]]!=alias[i-1]" by "ally[alias[i-1]]>=alias[i-1]" in line 553 in file "seq.c"
	$\mu_3$	Replace "i<b" by "i<=" in line 1097 in file "seq.c"
	$\mu_4$	Replace "j=i+1" by "j=i-1" in line 555 in file "seq.c"
	$\mu$	Replace "i<=(long)O" by "i>=(long)O" in line 994 in file "seq.c"
	$\mu$	Replace "itemp=alias[i-1]" by "itemp=alias[i+1]" in line 563 in file "seq.c"
	$\mu$	Replace "j<=(long)O" by "j>=(long)O" in line 1119 in file "seq.c"
	$\mu$	Replace "p->numsteps[i]+=weight[i]" by "p->numsteps[i]=weight[i]" in line 946 in file "seq.c"
	$\mu_{\infty}$	Replace "j<=(long)O" by "j>=(long)O" in line 1126 in file "seq.c"
	$\mu_{10}$	Replace "(i=)" by "(i!=)" in line 2700 in file "seq.c"
KNN	$\mu_1$	Replace "-" by "/" in line 29 in function "euclideanDistance"
	$\mu_2$	Replace "-" by "+" in line 29 in function "euclideanDistance"
	$\mu_3$	Replace "+=" by "=" in line 29 in function "euclideanDistance"
	$\mu_4$	Replace "math.sqrt(distance)" by "distance" in line 32 in function "euclideanDistance"
	$\mu$	Replace "reverse=True" by "reverse=False" in line 41 in function "getNeighbors"
	$\mu$	Swap lines 57 and 59 in function "getResponse"
	$\mu$	Replace "+=" by "*=" in line 52 in function "getResponse"
$\mu$	Replace "+=" by "=" in line 52 in function "getResponse"	
$\mu_{\infty}$	Replace "reverse=True" by "reverse=False" in line 55 in function "getResponse"	

- (a) both  $MR$  and  $MR$  belong to the special class of MRs in accordance with Definition 1 ;
  - (b)  $MR$  is composable with  $MR$  according to Definition 2 ;
  - (c)  $I(T) = T$ ,  $I$  is bijective, and  $O$  is injective.

In the above guideline, condition (c) involves applying Theorems 2, 3, and 4 as discussed in Section 4. This condition also involves the hypothesis stated at the beginning of Section 4.4, which was confirmed by our empirical analysis (Section 5) to be highly likely to be held true. When the injectivity/bijection requirement of our guideline does not hold, testers should consider other information about the program under test, and the amount of testing resources available, to inform their own decisions on MR composition. For example, if the program has a long execution time, then the reduction in program executions achieved by using composite MRs (even at the expense of a slight deterioration in fault detection effectiveness) may still be a better choice.

Table 1 summarizes the fault detection capability of  $MR$  in four different testing scenarios. It can be seen from Table 1 that, in three of the four scenarios (Scenarios 1, 2, and 3), the fault detection capability of the composite MR (i.e.,  $MR$ ) is identical to those of applying both  $MR$  and  $MR$ , if the three preconditions of the general guideline are satisfied. Since fewer test cases are required for testing  $MR$  when compared with using both  $MR$  and  $MR$ , the cost-effectiveness of MT in using  $MR$  (instead of  $MR$

and  $MR$ ) is obviously improved in these three scenarios. Even in Scenario 4 where a definite conclusion on the fault detection capability of  $MR$  (when compared with  $MR$  and  $MR$ ) cannot be drawn, we argue that it is very likely that the cost-effectiveness of  $MR$  is higher than  $MR$  and  $MR$  because of two reasons: (a) the situation where  $\theta = 0$  should rarely occur according to our theoretical analysis (see Section 4.4), which was further confirmed to be true by our empirical analysis (see observation 1 in Section 5.5); and (b) our empirical analysis showed that  $\theta \leq \min\{\theta, \theta\}$  for all cases and, in about three-quarter of the cases, we have  $\theta \leq \max\{\theta, \theta\}$  (see observations 2 and 3 in Section 5.5).

## 6.2 Applicability of our General Guideline

To evaluate the practicality and usefulness of our general guideline on MR composition, we reviewed a set of published papers on MT through which the following two questions could be answered:

- Q1: How likely is it that a given MR belongs to the special class, in accordance with Definition 1?
- Q2: Given an MR that belongs to the special class, according to Definition 1, how likely is this MR to have a bijective input mapping  $I$  and an injective output mapping  $O$ ?

We first studied in detail two recent survey papers on MT [37], [38], and then identified some other works on MT that have been developed after publishing those two surveys. After these exercises, we identified 10 popular application domains for MT as shown in Table 6. For each of these domains, we found some relevant published works related to MT. After a close examination, we compiled a list of MRs which were mentioned in these published works. Further checking of the list allowed us to identify some "common" MRs with similar types or characteristics within the same domain and even across different domains. These "common" MRs were counted only once in our review. For example, among the three MT-related papers on compilers, after tallying the count for "common" MRs, we found eight distinct MRs. We caution readers that, although our review did not (which was also infeasible to) involve every MT-related paper, we argue that the compiled list of MRs from our collected papers was fairly comprehensive because the list of MRs covered 10 different and popular domains for MT.

For Q1, we found 54.88% of MRs belong to the special class, in accordance with Definition 1. For Q2, we found that, among those MRs complying with Definition 1, 91.11% of them have a bijective  $I$  and an injective  $O$ . Based on these findings for Q1 and Q2, we can conclude that the three preconditions in our general MR composition guideline can easily be met (50.00% 54.88% 91.11%). This shows that our general guideline should be widely applicable to many testing scenarios and application domains.

We also noted that the applicability of our guideline varies across different application domains. More specifically, according to the results in 6, the guideline is largely applicable to compilers, numeric and scientific programs, and AI systems (e.g., image processing and autonomous car systems), and relatively less applicable to biomedical applications, web services, embedded systems, and online



TABLE 6: Applicability of the Guideline

Application Domains	References	No. of Identified MRs	No. of MRs complying with Def. 1	No. of MRs whose $I$ is bijective and $O$ is injective	Answer to Q1 (%)	Answer to Q2 (%)
Biomedical applications	[6], [7], [8]	42	12	12	28.57	100.00
Web services	[9], [10]	9	3	3	33.33	100.00
Embedded systems	[11], [12]	3	0	0	0.00	0.00
Component-based software	[13]	3	2	2	66.67	100.00
Compilers	[14], [15], [16]	8	8	5	100.00	62.50
Machine learning classifiers	[17], [18], [19], [20], [21]	16	7	7	43.75	100.00
Online search engines	[22], [23], [24]	9	0	0	0.00	0.00
Assorted computer science algorithms	[33], [34], [35]	27	16	11	59.26	68.75
Numerical and scientific programs	[30], [36]	17	14	14	82.35	100.00
AI systems (e.g., image processing and autonomous car systems)	[25], [26], [27], [28]	30	28	28	93.33	100.00
<b>Total</b>		<b>164</b>	<b>90</b>	<b>82</b>	<b>54.88</b>	<b>91.11</b>

### 6.3 Related Work

Two major challenges for MT are: (a) identification of MRs; and (b) additional computations of the program executions for follow-up test cases. Although MR composition can address both challenges, only few papers primarily focused on the effectiveness of MR composition — we found only two papers [30], [31] in this area. Obviously, MR composition reduces the number of program executions and, hence, lowers the computation costs — this is indisputable. However, these two studies on MR composition [30], [31] do not have a consensus on whether or not the fault detection capability after composition will be jeopardized. Furthermore, both studies [30], [31] adopted a purely empirical approach. Therefore, their observations could not provide a full picture of this issue and are dependent on the subjects being investigated. Understandably, some of their observations may look contradictory, that is, they reach different conclusions on the fault detection effectiveness of the composite MRs. On the other hand, with our theoretical results, such illusive contradictions are clarified. In summary, their results [30], [31] motivated our study, which in turn provides a more comprehensive interpretation of their results.

MR composition is an obvious and straightforward method to generate new MRs that is easily implemented and automated. However, this method requires the existence of some MRs for generation of new ones. Recently, with the increasing recognition and acceptance of MT by the software testing community, a growing number of research studies on MR generation/identification has emerged. Examples of these studies include machine-learning-based techniques [39], [40], [41], search-based techniques [42], [43], data-mutation-based techniques [26], [44], [45], pattern-based techniques [46], [47], and the category-choice approach [48], [49]. Since the main focus of this paper is not on MR generation, comparing our work with the above studies is beyond the scope of this paper.

## 7 SUMMARY AND CONCLUSION

Two major advantages of MR composition are the facilitation of automatic MR generation, and the reduction in testing costs (by reducing the number of program executions for follow-up test cases). However, MR composition has a

potential drawback: The fault detection capability of the composite MR may be lower than that of its component MRs, jeopardizing the overall effectiveness of MT. This issue motivated us to perform theoretical and empirical analyses, with a goal of identifying characteristics that the component MRs should possess so that the fault detection capability of the generated composite MR will likely not be less than that of its component MRs. In short, given a pair of metamorphic relations  $MR$  and  $MR$  that belong to the special class defined in Definition 1, where  $MR$  is composable with  $MR$ , they should be used to form a composite  $MR$  if both the following conditions are met: (a) the output mapping of  $MR$  is injective; and (b) the input mapping of  $MR$  is a bijective mapping from the source inputs of  $MR$  to the source inputs of  $MR$ . This result is produced based on Theorems 1 to 4, Propositions 1 and 2, and the empirical analysis discussed in the paper. This result provides a solid foundation for MT and has paved a path for future studies on MR composition.

Based on our theoretical and empirical analyses, a convenient yet effective general guideline on MR composition has been developed. We further performed studies (by using a sample of MRs extracted from previously published works on MT) to confirm the applicability of the guideline across a range of application domains.

Our study has focused on a special class of MRs, as stated in Definition 1. Although this class of MRs is common, it would be worthwhile to extend our study to examine MRs that are outside this class. A hierarchy of composable MRs and their composite MRs could then be built. It would be interesting to investigate what further information and insights could be exploited from this hierarchy, using the theoretical results reported in this paper. This future study would enhance the foundation of MT research. Another potentially fruitful direction is a large-scale empirical analysis of the relationships among  $\theta$ ,  $\min f\theta$ ,  $\theta g$ , and  $\max f\theta$ ,  $\theta g$ . The results of this empirical analysis would help practitioners better estimate their testing costs.

## ACKNOWLEDGMENTS

We are indebted to Dr. Dave Towey of the University of Nottingham Ningbo China for his valuable comments and

suggestions on improving this paper. This work was supported in part by the National Natural Science Foundation of China (Grant Nos. 61772055 and 61872169), the Technical Foundation Project of Ministry of Industry and Information Technology of China (Grant No. JSZL2016601B003), and the Equipment Preliminary R&D Project of China (Grant No. 41402020102).

## REFERENCES

- [1] P. Ammann and J. Offutt, *Introduction to Software Testing*. New York, NY: Cambridge University Press, 2018.
- [2] M. Pezzè and M. Young, *Software Testing and Analysis: Process, Principles, and Techniques*. Noida, India: Wiley, 2007.
- [3] E. T. Barr, M. Harman, P. McMinn, M. Shahbaz, and S. Yoo, "The oracle problem in software testing: A survey," *IEEE Transactions on Software Engineering*, vol. 41, no. 5, pp. 507–525, 2015.
- [4] T. Y. Chen, S. C. Cheung, and S. M. Yiu, "Metamorphic testing: A new approach for generating next test cases," Hong Kong University of Science and Technology, Hong Kong, Technical Report HKUST-CS98-01, 1998.
- [5] F. T. Chan, T. Y. Chen, S. C. Cheung, M. F. Lau, and S. M. Yiu, "Application of metamorphic testing in numerical analysis," in *Proceedings of the IASTED International Conference on Software Engineering (SE'98)*, 1998, pp. 191–197.
- [6] T. Y. Chen, J. W. K. Ho, H. Liu, and X. Xie, "An innovative approach for testing bioinformatics programs using metamorphic testing," *BMC Bioinformatics*, vol. 10, no. 1, 2009, Article no. 24.
- [7] L. L. Pullum and O. Ozmen, "Early results from metamorphic testing of epidemiological models," in *Proceedings of ASE/IEEE International Conference on BioMedical Computing (BioMedCom)*, 2012, pp. 62–67.
- [8] M. S. Sadi, F.-C. Kuo, J. W. K. Ho, M. A. Charleston, and T. Y. Chen, "Verification of phylogenetic inference programs using metamorphic testing," *Journal of Bioinformatics and Computational Biology*, vol. 9, no. 6, pp. 729–747, 2011.
- [9] W. K. Chan, S. C. Cheung, and K. R. P. Leung, "A metamorphic testing approach for online testing of service-oriented software applications," *International Journal of Web Services Research (IJWSR)*, vol. 4, no. 2, pp. 61–81, 2007.
- [10] C.-A. Sun, G. Wang, B. Mu, H. Liu, Z. S. Wang, and T. Y. Chen, "Metamorphic testing for web services: Framework and a case study," in *Proceedings of 2011 IEEE International Conference on Web Services*, 2011, pp. 283–290.
- [11] W. K. Chan, T. Y. Chen, H. Lu, T. H. Tse, and S. S. Yau, "Integration testing of context-sensitive middleware-based applications: A metamorphic approach," *International Journal of Software Engineering and Knowledge Engineering*, vol. 16, no. 5, pp. 677–703, 2006.
- [12] F.-C. Kuo, T. Y. Chen, and W. K. Tam, "Testing embedded software by metamorphic testing: A wireless metering system case study," in *Proceedings of 36th Conference on Local Computer Networks*, 2011, pp. 291–294.
- [13] X.-L. Lu, Y.-W. Dong, and C. Luo, "Testing of component-based software: A metamorphic testing methodology," in *Proceedings of 7th International Conference on Ubiquitous Intelligence & Computing and 7th International Conference on Autonomic & Trusted Computing*, 2010, pp. 272–276.
- [14] V. Le, M. Afshari, and Z. Su, "Compiler validation via equivalence modulo inputs," *ACM SIGPLAN Notices*, vol. 49, no. 6, pp. 216–226, 2014.
- [15] Q. Tao, W. Wu, C. Zhao, and W. Shen, "An automatic testing approach for compiler based on metamorphic testing technique," in *Proceedings of 2010 Asia-Pacific Software Engineering Conference*, 2010, pp. 270–279.
- [16] A. F. Donaldson, H. Evrard, A. Lascu, and P. Thomson, "Automated testing of graphics shader compilers," in *Proceedings of ACM on Programming Languages*, 2017, Article no. 93.
- [17] C. Murphy, G. E. Kaiser, and L. Hu, "Properties of machine learning applications for use in metamorphic testing," Columbia University, U.S., Technical Report CUCS-011-08, 2011.
- [18] X. Xie, J. Ho, C. Murphy, G. Kaiser, B. Xu, and T. Y. Chen, "Application of metamorphic testing to supervised classifiers," in *Proceedings of 9th International Conference on Quality Software*, 2009, pp. 135–144.
- [19] X. Xie, J. W. K. Ho, C. Murphy, G. Kaiser, B. Xu, and T. Y. Chen, "Testing and validating machine learning classifiers by metamorphic testing," *Journal of Systems and Software*, vol. 84, no. 4, pp. 544–558, 2011.
- [20] S. Nakajima and H. N. Bui, "Dataset coverage for testing machine learning computer programs," in *Proceedings of 23rd Asia-Pacific Software Engineering Conference*, 2016, pp. 297–304.
- [21] P. Saha and U. Kanewala, "Fault detection effectiveness of metamorphic relations developed for testing supervised classifiers," in *Proceedings of IEEE International Conference on Artificial Intelligence Testing (AITest)*, 2019, pp. 157–164.
- [22] Z. Q. Zhou, S. Zhang, M. Hagenbuchner, T. H. Tse, F.-C. Kuo, and T. Y. Chen, "Automated functional testing of online search services," *Software Testing, Verification and Reliability*, vol. 22, no. 4, pp. 221–243, 2012.
- [23] Z. Q. Zhou, S. Xiang, and T. Y. Chen, "Metamorphic testing for software quality assessment: A study of search engines," *IEEE Transactions on Software Engineering*, vol. 42, no. 3, pp. 264–284, 2015.
- [24] J. Brown, Z. Q. Zhou, and Y.-W. Chow, "Metamorphic testing of navigation software: A pilot study with Google maps," in *Proceedings of 51st Hawaii International Conference on System Sciences*, 2018, pp. 5687–5696.
- [25] J. Mayer and R. Guderlei, "On random testing of image processing applications," in *Proceedings of 6th International Conference on Quality Software*, 2006, pp. 85–92.
- [26] H. Zhu, D. Liu, I. Bayley, R. Harrison, and F. Cuzzolin, "Data-morphic testing: A method for testing intelligent applications," in *Proceedings of IEEE International Conference on Artificial Intelligence Testing (AITest)*, 2019, pp. 149–156.
- [27] M. Zhang, Y. Zhang, L. Zhang, C. Liu, and S. Khurshid, "DeepRoad: GAN-based metamorphic testing and input validation framework for autonomous driving systems," in *Proceedings of 33rd ACM/IEEE International Conference on Automated Software Engineering*, 2018, pp. 132–142.
- [28] Y. Tian, K. Pei, S. Jana, and B. Ray, "DeepTest: Automated testing of deep-neural-network-driven autonomous cars," in *Proceedings of 40th International Conference on Software Engineering*, 2018, pp. 303–314.
- [29] Z. Q. Zhou and L. Sun, "Metamorphic testing of driverless cars," *Communications of the ACM*, vol. 62, no. 3, pp. 61–67, 2019.
- [30] G. Dong, B. Xu, L. Chen, C. Nie, and L. Wang, "Case studies on testing with compositional metamorphic relations," *Journal of Southeast University (English Edition)*, vol. 24, no. 4, pp. 437–443, 2008.
- [31] H. Liu, X. Liu, and T. Y. Chen, "A new method for constructing metamorphic relations," in *Proceedings of 12th International Conference on Quality Software*, 2012, pp. 59–68.
- [32] B. H. Smith and L. Williams, "On guiding the augmentation of an automated test suite via mutation analysis," *Empirical Software Engineering*, vol. 14, no. 3, pp. 341–369, 2009.
- [33] M. Jiang, T. Y. Chen, F.-C. Kuo, and Z. Ding, "Testing central processing unit scheduling algorithms using metamorphic testing," in *Proceedings of 4th International Conference on Software Engineering and Service Science*, 2013, pp. 530–536.
- [34] P. Rao, Z. Zheng, T. Y. Chen, N. Wang, and K. Cai, "Impacts of test suite's class imbalance on spectrum-based fault localization techniques," in *Proceedings of 13th International Conference on Quality Software*, 2013, pp. 260–267.
- [35] X. Xie, W. E. Wong, T. Y. Chen, and B. Xu, "Metamorphic slice: An application in spectrum-based fault localization," *Information and Software Technology*, vol. 55, no. 5, pp. 866–879, 2013.
- [36] T. Y. Chen, J. Feng, and T. H. Tse, "Metamorphic testing of programs on partial differential equations: A case study," in *Proceedings of 26th Annual International Computer Software and Applications*, 2002, pp. 327–333.
- [37] S. Segura, G. Fraser, A. B. Sanchez, and A. Ruiz-Cortés, "A survey on metamorphic testing," *IEEE Transactions on Software Engineering*, vol. 42, no. 9, pp. 805–824, 2016.
- [38] T. Y. Chen, F.-C. Kuo, H. Liu, P.-L. Poon, D. Towey, T. H. Tse, and Z. Q. Zhou, "Metamorphic testing: A review of challenges and opportunities," *ACM Computing Surveys*, vol. 51, no. 1, article 4, pp. 1–27, 2018.
- [39] U. Kanewala and J. M. Bieman, "Using machine learning techniques to detect metamorphic relations for programs without test oracles," in *Proceedings of 2013 IEEE 24th International Symposium on Software Reliability Engineering (ISSRE)*, 2013, pp. 1–10.



- [40] U. Kanewala, J. M. Bieman, and A. Ben-Hur, "Predicting metamorphic relations for testing scientific software: A machine learning approach using graph kernels," *Software Testing, Verification and Reliability*, vol. 26, no. 3, pp. 245–269, 2016.
- [41] A. Nair, K. Meinke, and S. Eldh, "Leveraging mutants for automatic prediction of metamorphic relations using machine learning," in *Proceedings of 3rd ACM SIGSOFT International Workshop on Machine Learning Techniques for Software Quality Evaluation*, 2019, pp. 1–6.
- [42] J. Zhang, J. Chen, D. Hao, Y. Xiong, B. Xie, L. Zhang, and H. Mei, "Search-based inference of polynomial metamorphic relations," in *Proceedings of 29th ACM/IEEE International Conference on Automated Software Engineering*, 2014, pp. 701–712.
- [43] B. Zhang, H. Zhang, J. Chen, D. Hao, and P. Moscato, "Automatic discovery and cleansing of numerical metamorphic relations," in *Proceedings of IEEE International Conference on Software Maintenance and Evolution (ICSME)*, 2019, pp. 235–245.
- [44] H. Zhu, "JFuzz: A tool for automated Java unit testing based on data mutation and metamorphic testing methods," in *Proceedings of 2nd International Conference on Trustworthy Systems and Their Applications*, 2015, pp. 8–15.
- [45] C.-A. Sun, Y. Liu, Z. Wang, and W. K. Chan, " $\mu$ MT: A data mutation directed metamorphic relation acquisition methodology," in *Proceedings of 2016 IEEE/ACM 1st International Workshop on Metamorphic Testing (MET)*, 2016, pp. 12–18.
- [46] S. Segura, J. A. Parejo, J. Troya, and A. Ruiz-Cortés, "Metamorphic testing of RESTful web APIs," *IEEE Transactions on Software Engineering*, vol. 44, no. 11, pp. 1083–1099, 2017.
- [47] Z. Q. Zhou, L. Sun, T. Y. Chen, and D. Towey, "Metamorphic relations for enhancing system understanding and use," *IEEE Transactions on Software Engineering*, 2018. doi: 10.1109/TSE.2018.2876433.
- [48] T. Y. Chen, P.-L. Poon, and X. Xie, "METRIC: Metamorphic relation identification based on the category-choice framework," *Journal of Systems and Software*, vol. 116, pp. 177–190, 2016.
- [49] C.-A. Sun, A. Fu, P.-L. Poon, X. Xie, H. Liu, and T. Y. Chen, "METRIC+: A metamorphic relation identification technique based on input plus output domains," *IEEE Transactions on Software Engineering* (in press).



**Pak-Lok Poon (M' 01)** received his Ph.D degree in software engineering from The University of Melbourne. He is an Associate Professor with the School of Engineering and Technology, Central Queensland University, Australia. His research interests include software testing, requirements engineering and inspection, electronic commerce, and computers in education. He was a guest editor of the special issue of the *Journal of Systems and Software* on test oracles in 2018, and will be a guest editor of the special issue of the same journal on metamorphic testing in 2021. He was also an organizer for the *3rd, 4th, and 5th International Workshops on Metamorphic Testing* in 2018, 2019, and 2020, respectively.



**Kun Qiu** received his B.S. in Automation from the Hefei University of Technology, Hefei, China, in 2013. He is currently a Ph.D candidate in the School of Automation Science and Electrical Engineering at Beihang University, Beijing, China. His research interests include software testing and software reliability analysis.



**Zheng Zheng (SM' 18)** received his Ph.D degree in computer software and theory from the Chinese Academy of Science, Beijing, China. In 2014, he worked as a research scholar in the Department of Electrical and Computer Engineering at Duke University, Durham, North Carolina, USA. He is currently a Professor at Beihang University, Beijing, China. His research interests include software dependability modeling, software testing, and software fault localization.



**Tsong Yueh Chen (M' 03)** received his Ph.D degree from The University of Melbourne. He is currently a Professor of Software Engineering at Swinburne University of Technology, Australia. Prior to joining Swinburne, he taught at The University of Hong Kong and The University of Melbourne. He is the inventor of metamorphic testing and adaptive random testing.